System-Wide Water SWWRP Resources Program

Performance Study of

Parallel Algorithms in DWASH123D





Background

Based on the first-principle, physics-based numerical model WASH120pt, ph. U.S.
Army Engineer Research and Developing, a personal research and the properties of the propertie

combination of one-dimensional (1-D) channel network, 2-D overfand regimes, and
3-0 subsurface modals. This poster presents the outcome of a performance study on
the parallel algorithms cerrently employed in pWASH123D. The experimental area
includes a 576-square-mile domain covering most of the land south of the Tallar
Trail in South Florida and north of the Gulf of Mexico, Florida Bay, and Biscayne
Bay, This area was discretized to there mesh resolutions—and sicretized to there mesh resolutions—of
since the control of the control of

pWASH123D

- Solve the cross-sectional area-averaged diffusive wave equation with the semi-Lagrangian finite element method (FEM) for 1-D channel flow.
 Solve the depth-averaged diffusive wave equation with the semi-Lagrangian FEM for
- 2-D overland flow.
- Solve the Richards' equation for variably saturated porous media with the Galerkin FEM for 3-D subsurface flow.
- Enforce the continuity of flux and/or state variables (e.g., water head, concentration) on the interface of two media.
- Consider rainfall, evapotranspiration, rule-controlled flow, injection/withdrawal as sources/sinks at respective media.
- sourcesisms at respective media.

 Employ different time intervals for computations in different dimensions to resolve various flow processes $(\Delta t_{-} \ge \Delta t_{-} = 0.5 \text{ s. } \Delta t_{-} = 30 \text{ m. } \Delta t_{-} = 5 \text{ s. } \Delta t_{-} = 0.5 \text{ s. } \text{ used in }$
- Implement independent domain partitioning, coherent data structures, and interactions between domains.





Hierarchical Data Structures



Coupling Algorithm



Parallelization Strategy

- \bullet DBuilder is used in both 2- and 3-D computation to balance computational load on each processor.
- Each processor reads complete 1-D channel information and executes 1-D computation without partitioning to avoid excessive run time overhead from data exchange among processors.

Test Example



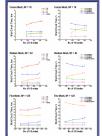


POCs:
Jing-Ru (Ruth) C. Cheng
Phone: (601) 634-4692 • E-Mailt: ruth.c.nbeng@erdc.usace.army.mil
Hwai-Ping (Pearce) Cheng, Hsin-Chi (Jerry) Lin,
Robert R. Richards. David R. Richards. Earl V. Edris

https://swwrp.usace.army.mil

Computational Meshes

NP	16		64			
					128	
Components	ND	NE	ND	NE	ND	NE
2-0	8,487	16,583	42,941	84,996	101,148	200,935
3-0	59,409	99,498	558,233	1,019,952	1,314,924	2,411,220
1-0	ND	NR.	ND	NR.	ND	NR.
Case 1	89	- 1	206	- 1	306	- 1
Case 2	127	2	298	2	492	2
Case 3	200	4	463	4	710	4
Case 4	214	5	495	5	760	5



Results

 The wall-clock time for 1-D computation is approximately proportional to the number of 1-D nodes included in the simulation, which results in the increase of both total wall-clock time and the wallclock time percent for 1-D computation.

- The wall-clock time for 2-and 3-D computations basically does not vary with the number of 1-D nodes. With the current parallelization strategy in pWASH123D, it is obvious that the fewer 1-D nodes considered for computation, the less time spent for 1-D computation.
- The time spent in couplers is negligible when compared with that spent in 1-, 2-, or 3-D computation.

Summary and Future Plans

The 1-D computation significantly takes up a portion of overall wall-tock time of simulations when using a larger number of processors with the current parallelization strategy. This result is also strongly related to the multiple time-step coupling algorithm employed to resolve various physical processes. An autonomous approach, which can guarantee convergence of the nonlinear system using larger time-step sizes, will highly benefit such a parallel watershed model. It is also worthwhile to investigate time-scape can arallelism on the lower dimensional domains.